

Introduction to Language Theory and Compilation Solutions

Session 6: First sets, Follow sets and LL(1) parsing

Solutions

Ex. 1. With regards to the grammar:

1. Give the $\text{First}^1(A)$ and the $\text{Follow}^1(A)$ sets for each $A \in V$.

Symbol	$\text{First}^1()$	$\text{Follow}^1()$
S	begin	
program	begin	\$
statement list	ID read write	end
statement tail	ID read write ϵ	end
statement	ID read write	ID read write end
id list	ID	)
id tail	, ϵ	)
expr list	(ID INTLIT	)
expr tail	, ϵ	)
expression	(ID INTLIT	; ,)
primary tail	+ - ϵ	; ,)
primary	(ID INTLIT	+ - ; ,)
add op	+ -	(ID INTLIT

2. Give the $\text{First}^2(<\text{expression}>)$ and the $\text{Follow}^2(<\text{expression}>)$ sets.

$\text{First}^2(<\text{expression}>) = \{ (, (ID, (INTLIT, ID+, ID-, INTLIT+, INTLIT-, ID, INTLIT) \}$

$\text{Follow}^2(<\text{expression}>) = \{ , (, , ID, , INTLIT,); ,)+,)-,), ,), , ;read, ;write, ;ID, ;end \}$

Ex. 2. Which grammars are LL(1)?

1. not LL(1) because $a \in \text{Follow}^1(A)$, $a \in \text{First}^1(A)$ and we have the rule $A \rightarrow \epsilon$. Thus, $M[A, a] = \{\text{Produce 2, Produce 3}\}$.
2. LL(1).
3. LL(1).
4. not LL(1) because $b \in \text{First}^1(A)$, $b \in \text{Follow}^1(B)$ and we have the rule $B \rightarrow \epsilon$. Thus, $M[A, b] = \{\text{Produce 3, Produce 4}\}$.

Ex. 3. Give the action table of the grammar:

	-	(	)	ID	\$
<S>	1	1		1	
<Expr>	2	3		4	
<ExprTail>	5		6		6
<Var>				7	
<VarTail>	9	8	9		9

Ex. 4. Program a recursive descent parser (in C, C++, ...) for rules (14) through (21).

```

void void expression(){
    primary(); primary_tail();
}
void primary_tail(){
    token tok = next_token();
    // primary_tail is nullable, we check with the follow set
    // and return if primary_tail has to be epsilon
    switch(tok){
        case ' ';:
        case ',':
        case '(': return;
    }
    // otherwise, we try to match the rule
    add_op(); primary(); primary_tail();
}
void primary(){
    token tok = next_token();
    switch(tok){
        case '(':
            match('('); expression(); match(')');
            break;
        case ID:
            match(ID);
            break;
        case INTLIT:
            match(INTLIT);
            break;
        default:
            syntax_error(tok); break;
    }
}
void add_op(){
    token tok = next_token();
    switch(tok){
        case '+':
            match('+');
            break;
        case '-':
            match('-');
            break;
        default:
            syntax_error(tok); break;
    }
}

```